

# Midterm Review

---

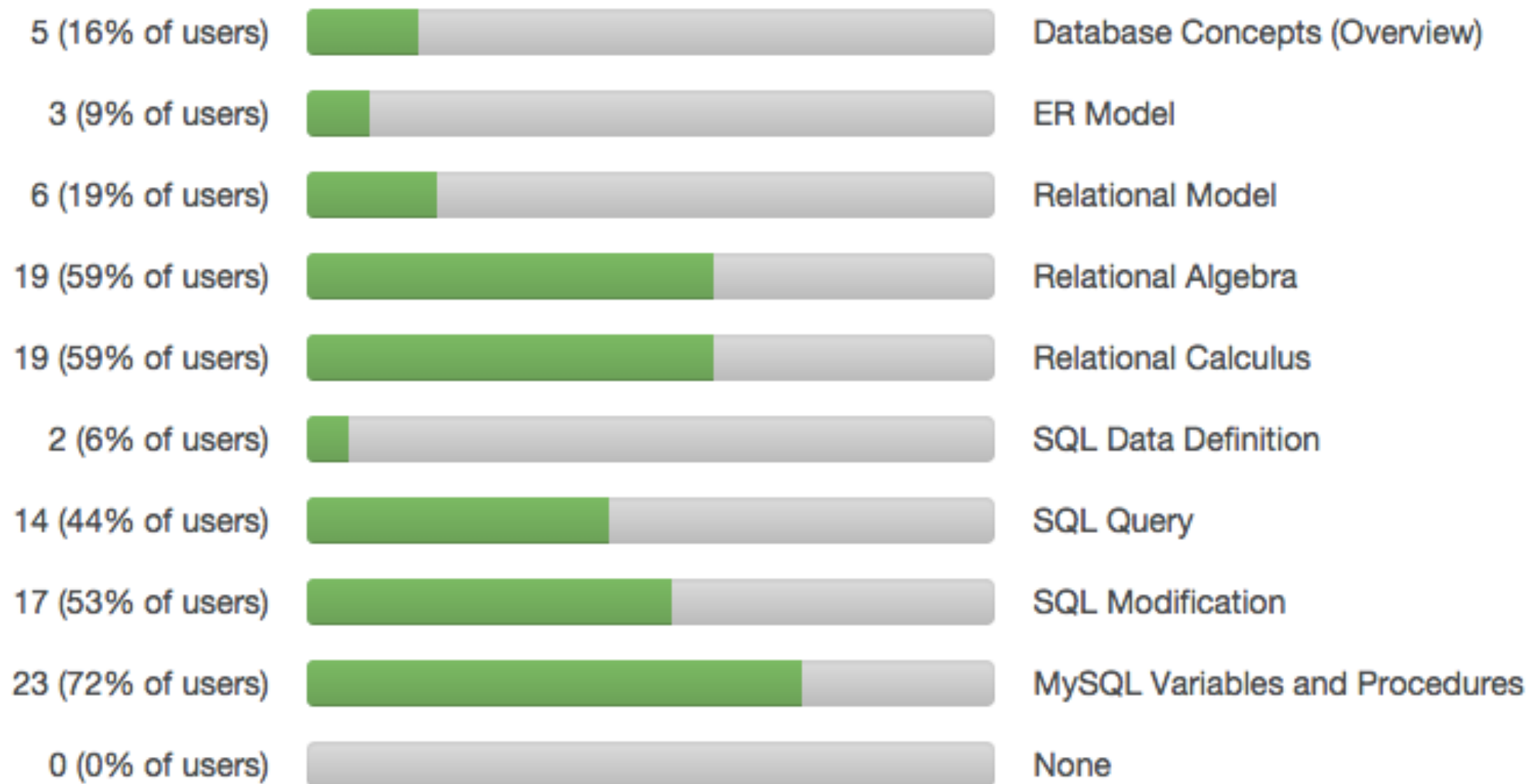
CS 377: Database Systems

# Piazza Poll Results

---

## Midterm Review Topics closes in 2 day(s)

A total of 32 vote(s) in 104 hours



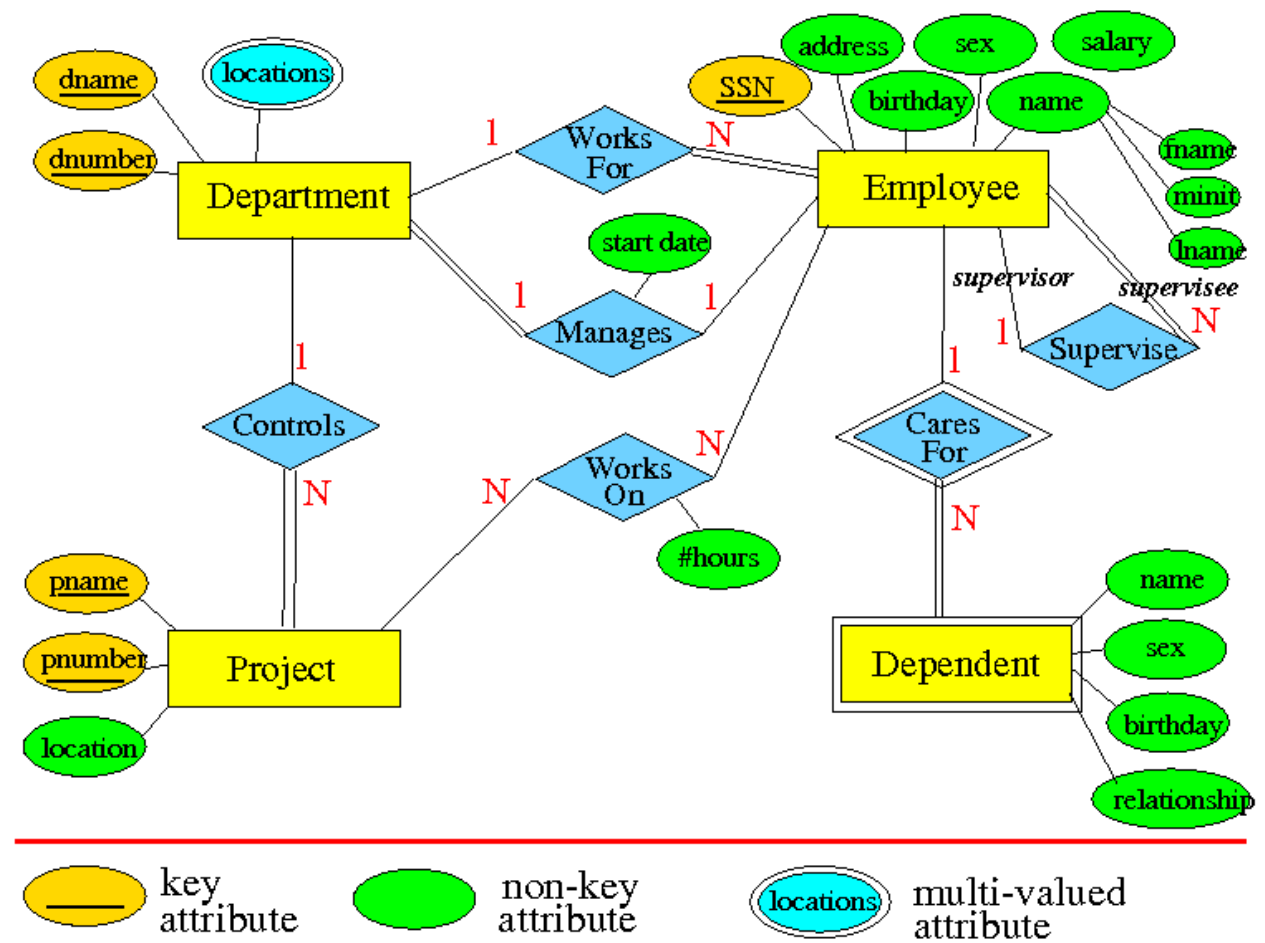
# Database Concepts

---

- Data model categories: high-level or conceptual data models, low-level or physical data models, and representational or implementation data models
- Physical data and logical data independence
  - How metadata fits into the picture
- Three schema architecture

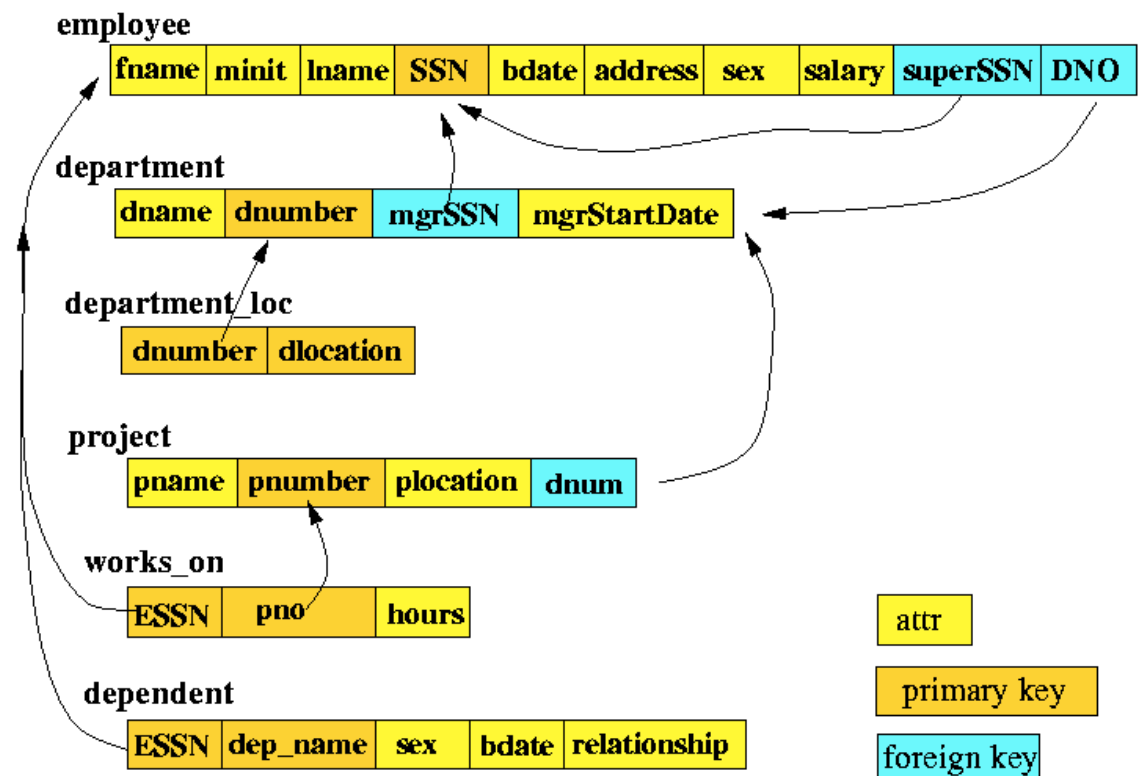
# Entity Relationship (ER) Model

- Entity
  - Attributes
    - Weak Entity
- Relationship
  - Degree
  - Cardinality ratio constraint
  - Participation constraint



# Relation Model

- Relation, attributes
- Schema vs instance
- Relational model constraints
  - Domain constraint
  - Key constraint
  - Referential integrity constraint



# ER to Relational Model

---

Entity set and relationships and convert them to relation

ER Model	Relational model
Entity type	Entity relation
1:1 or 1:N relationship	Expand (or create R relation)
M:N relationship	Create R relation with two foreign keys
n-ary relationship type	Create R relation with n foreign keys
Simple attribute	Attribute
Composite attribute	Set of simple component attributes
Multivalued attribute	Relation and foreign key
Key attribute	Primary (or secondary) key

# Relational Algebra

Operation	Notation	Purpose
SELECT	$\sigma_{\langle \text{selection condition} \rangle}(R)$	Selects all tuples that satisfy the selection condition from a relation R
PROJECT	$\pi_{\langle \text{attribute list} \rangle}(R)$	New relation with subset of attributes of R and removes duplicate tuples
THETA_JOIN	$R_1 \bowtie_{\langle \text{join condition} \rangle} R_2$	All combinations of tuples from R <sub>1</sub> and R <sub>2</sub> that satisfy the join condition
EQUIJOIN	$R_1 \bowtie_{\langle \text{join condition} \rangle} R_2$	Theta join with only equality join comparisons
NATURAL JOIN	$R_1 *_{\langle \text{join condition} \rangle} R_2$	Equijoin except join attributes of R <sub>2</sub> are not included in the resulting relation
UNION	$R_1 \cup R_2$	Relation that includes all tuples in R <sub>1</sub> or R <sub>2</sub>
INTERSECTION	$R_1 \cap R_2$	Relation that includes all tuples in both R <sub>1</sub> and R <sub>2</sub>
DIFFERENCE	$R_1 - R_2$	Relation that includes all tuples in R <sub>1</sub> that are not in R <sub>2</sub>
CARTESIAN PRODUCT	$R_1 \times R_2$	Relation with attributes of R <sub>1</sub> and R <sub>2</sub> and includes tuples with all possible combinations of tuples of R <sub>1</sub> and R <sub>2</sub>
DIVISION	$R_1(Z) \div R_2(Y)$	Relation that includes all tuples t[X] in R <sub>1</sub> (Z) that appear in R <sub>1</sub> in combination with every tuple from R <sub>2</sub> (Y) where $Z = X \cup Y$
GROUP BY AGGREGATE	$\langle \text{group attrs} \rangle \mathcal{F}_{\langle \text{set funcs} \rangle}$	Relation that includes the grouping attributes and the set function values

# Banking Example

---

BRANCH

<u>bname</u>	bcity	assets
--------------	-------	--------

CUSTOMER

<u>cname</u>	street	bcity
--------------	--------	-------

ACCOUNT

<u>aID</u>	bname	balance
------------	-------	---------

LOAN

<u>loanID</u>	bname	amount
---------------	-------	--------

DEPOSITOR

<u>cname</u>	<u>aID</u>
--------------	------------

BORROWER

<u>cname</u>	<u>loanID</u>
--------------	---------------

- Find the names of all customers who have a loan and a savings account at the bank
- Find the names of all customers who have a loan at the Decatur branch but do not have a savings account at any branch of the bank
- Find all customers who have a savings account at all branches located in Atlanta city



# Relational Calculus (Tuple Relational Calculus)

---

Query of the form:  $\{t \mid \text{CONDITION}(t)\}$

- Conditions are formulas and are recursively defined
- Atomic formula (Relation( $t$ ),  $R.a \text{ op } S.b / \text{constant}$ )
- Special formula quantifiers
  - Universal quantifier  $(\forall t) (\text{Condition}(t))$
  - Existential quantifier  $(\exists t) (\text{Condition}(t))$

# Banking Example

---

BRANCH

<u>bname</u>	bcity	assets
--------------	-------	--------

CUSTOMER

<u>cname</u>	street	bcity
--------------	--------	-------

ACCOUNT

<u>aID</u>	bname	balance
------------	-------	---------

LOAN

<u>loanID</u>	bname	amount
---------------	-------	--------

DEPOSITOR

<u>cname</u>	<u>aID</u>
--------------	------------

BORROWER

<u>cname</u>	<u>loanID</u>
--------------	---------------

- Find loan number for each loan of an amount greater than \$1200
- Find the names of all customers who have a loan and a savings account at the bank
- Find all customers who have a savings account at all branches located in Atlanta city

# SQL Data Definition

---

- Create database
- Create table
  - Attribute datatypes and constraints
  - Key constraints (primary and foreign key)
  - Circular integrity constraints
- Alter tables
  - Add/remove attributes
  - Add/remove constraints
- Drop tables & databases

# SQL Query

---

- Basic SQL query

```
SELECT    [DISTINCT] <attribute list>  
FROM      <table list>  
[WHERE    <condition on the tables>]  
[GROUP BY <grouping attributes>]  
[HAVING   <group condition>]  
[ORDER BY <attribute list>  ASC | DESC]  
[LIMIT    <number of tuples>]
```

- Nested queries and temporal relations
- Advanced query formulations

# SQL Data Modification

---

Data modification does not return a result but changes the database

- INSERT (add new tuples)
  - Literal values (constant or known values):  
**INSERT INTO <table name>[(<attr names>)]  
VALUES (<list of values>);**
  - Result from a **SELECT** command:  
**INSERT INTO <table name>[(<attr names>)]  
(<SELECT subquery>)**

# SQL Data Modification (2)

---

Data modification does not return a result but changes the database

- DELETE (remove tuples)  
**DELETE FROM <relation> WHERE <condition>;**
- Tuples are deleted from only one table at a time unless **CASCADE** is specified on a referential integrity constraint
- UPDATE (change value(s) of existing tuples)

# SQL Data Modification (3)

---

Data modification does not return a result but changes the database

- UPDATE (change value(s) of existing tuples)

**UPDATE** <relation>

**SET** <list of attribute assignments>

**WHERE** <condition>;

- Modify/change certain attributes in certain tuples of a relation
- Only changes tuples that match the WHERE condition

# SQL Views

---

- View is a virtual table that does not exist in physical form
  - Allows ability to present information in different ways to different users
  - Can be used like an ordinary relation and simplifies complex queries
  - Limits data access to specific users (sensitive data can be hidden)
  - If conceptual schema changes, only the **SELECT** query needed to construct view needs to change
- Syntax:  
**CREATE VIEW <name> AS <query>;**



# MySQL Session Variables

---

- A session starts with a connection to the SQL server and ends when the connection is closed
- Session variables can be created anytime during a SQL session and exists for the remainder of the SQL session
- Always begins with the symbol “@”  
(e.g, @x, @count)
- Syntax:  
**SET <varName> = express;**  
**SELECT ... INTO @varname FROM ... WHERE ...;**

# MySQL Temporary Tables

---

- Temporary tables are used to store and process intermediate results:

**CREATE TEMPORARY TABLE**

...

- Same selection, update, and join capabilities in typical SQL tables
- Deleted when the current client session terminates
- Stored functions

# MySQL Stored Procedures

---

- Generalization of SQL by adding programming language-like structure to the SQL language
- Syntax:  
**DELIMITER <DL>**  
**CREATE PROCEDURE <procedure name> (parameters)**  
**BEGIN**  
    <statements of the procedure>  
**END <DL>**
- A stored procedure can only be used within the database where the stored procedure was defined

# MySQL Stored Procedures (2)

---

- Stored procedure can have local variables
  - **BEGIN** and **END** keywords defines the scopes of local variables
  - Inner levels can access variables from outer levels but not vice-versa
- Stored procedure can have parameters (similar to methods in programming language)
  - 3 modes to pass in parameters (IN, OUT, INOUT)

# MySQL Stored Procedures (3)

---

Control structures are available similar to traditional programming languages

- IF statements (conditional)
  - **IF <condition> THEN <command> END IF;**
  - **IF <condition> THEN <command1>  
ELSE <command2> END IF;**
- CASE statements (alternative conditional structure)  
**CASE <case expression>**
  - WHEN <expression1> THEN <command1>**
  - ...**
  - ELSE <commandN>****END CASE;**

# MySQL Stored Procedures (4)

---

Control structures are available similar to traditional programming languages

- LOOP statements (repeated execution)
  - **WHILE <condition> DO <commands>  
END WHILE;**
  - **REPEAT <commands> UNTIL <condition>  
END REPEAT;**
  - **<LoopLabel>:  
LOOP  
    <commands>  
    IF <condition1> THEN LEAVE <LoopLabel>;  
    IF <condition2> THEN ITERATE <LoopLabel>;  
END LOOP;**

# MySQL Stored Functions

---

User-defined functions

```
CREATE FUNCTION <function_name>(parameter)  
  RETURNS datatype  
  [NOT] DETERMINISTIC  
<statements>;
```

- Returns a single value (similar to aggregate functions)
- Meant to encapsulate common formulas or business rules that are reusable
- Can be used in SQL SELECT statements

# JDBC Program Steps

---

- Import JDBC library (java.sql.\*)
- Load appropriate JDBC driver
- Create a connection object
- Create a statement object
- Submit SQL statement
- Process query results
- Close connections